# Laplacian Spectra of Graphs and Cyber-Insurance Protection

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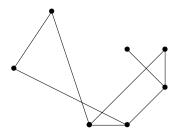
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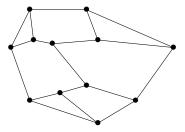
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- ► Or... Interconnected data, in the form of a graph → susceptible to data breaches, which can be insured.
- Question: impact of the topology of the network on the chosen protection?

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- Nodes can represent: computers, servers, softwares, routers, data sets, etc.
- Edges can be: internet connections, distant access to servers, links between softwares etc.



(a) A graph on n = 7 vertices, with 8 edges.



(b) A 3-regular graph on n = 12 vertices.

#### Outline of the talk

- ► The SIS model.
- An optimization criteria: the **algebraic connectivity**.
- Solution and numerical illustrations.

#### Literature

- Economics of information security (Gordon and Loeb 2002, Böhme 2013).
- Insurability of cyber risk (Eling 2020's)
- Cyber claims pricing (Fahrenwaldt & al 2018, Xu and Hua 2019)
- Impact of the network's topology (Hillairet & al 2021, Hillairet and Lopez 2021)
- Statistical inference (Bessy-Rolland & al 2020, Farkas & al 2021)
- And in practice ? (Romanosky &al 2019, Malavasi & al 2021)
- Eigenvalues optimization, epidemic models (refs. in the paper).

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# Inhomogeneous Suseptible-Infected-Susceptible (SIS) model on Edges.

State at time t of edge i is Bernoulli: X<sub>i</sub>(t) = 0 when edge i is healthy and X<sub>i</sub>(t) = 1 when edge i is infected.

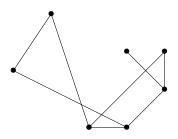
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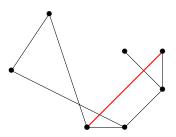
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- When infected, edge *i* can infect each edge *j* sharing a common vertex with rate β<sub>ij</sub> > 0.
- Each edge *i* can be cured, i.e. its state  $X_i$  jumps from 1 to 0 at rate  $\delta_i > 0$ .
- ▶ The Poisson infection and curing processes are assumed independent.

Inhomogeneous Suseptible-Infected-Susceptible (SIS) model.



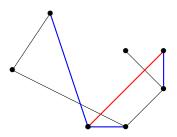
No edge is infected.

Inhomogeneous Suseptible-Infected-Susceptible (SIS) model.



Red edge *i* is infected, and can be cured at rate  $\delta_i > 0$ .

Inhomogeneous Suseptible-Infected-Susceptible (SIS) model.



Edge *i* infects each blue edge *j* at rate  $\beta_{ij} > 0$ .

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- Long time steady state: which is accurate when infection probabilities are high, corresponding to cascading infections.

Let  $\tilde{a} =$  edge adjacency matrix and m = number of edges.

$$rac{X_i(t+\Delta t)-X_i(t)}{\Delta t}=(1-X_i(t))\sum_{j=1}^m ilde{a}_{ij}eta_{ij}X_j(t)-\delta_iX_i(t).$$

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leads to

$$rac{d v_i(t)}{dt} = \sum_{j=1}^m ilde{a}_{ij} eta_{ij} v_j(t) - \sum_{j=1}^m ilde{a}_{ij} eta_{ij} \mathbb{E}[X_i(t)X_j(t)] - \delta_i v_i(t),$$

where  $v_i(t) = \mathbb{P}(X_i(t) = 1)$ .

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- System of *m* coupled ODEs.
- System of  $m 2^m 1$  coupled ODEs.
- Curse of dimensionality!

N-Intertwined Mean Field Approximation:

$$rac{d \mathsf{v}_i(t)}{dt} = \sum_{j=1}^m ilde{a}_{ij} eta_{ij} \mathsf{v}_j(t) - \sum_{j=1}^m ilde{a}_{ij} eta_{ij} \mathbb{E}[\mathsf{X}_i(t)] \mathbb{E}[\mathsf{X}_j(t)] - \delta_i \mathsf{v}_i(t),$$

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i.e

$$\frac{d\mathsf{v}_i(t)}{dt} = \left(\sum_{j=1}^m \tilde{\mathsf{a}}_{ij}\beta_{ij}\mathsf{v}_j(t)\right)(1-\mathsf{v}_i(t)) - \delta_i\mathsf{v}_i(t), \quad i=1,\ldots,m.$$

*N*-Intertwined Mean Field Approximation and Long time steady state:

$$\frac{dv_i(t)}{dt} = 0 = \left(\sum_{j=1}^m \tilde{a}_{ij}\beta_{ij}v_j(t)\right)(1-v_i(t)) - \delta_i v_i(t), \quad i = 1, \dots, m.$$

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i.e.

$$v_i = \frac{\beta \sum_{j=1}^m \tilde{a}_{ij} v_j}{\delta_i + \beta \sum_{j=1}^m \tilde{a}_{ij} v_j}, \quad i = 1, \dots m_i$$

where  $\beta_{ij} = \beta$ .

# **Optimization Problem**

Control the vector of curing rates (δ<sub>1</sub>,..., δ<sub>m</sub>) to make the graph as connected as possible, in a situation of cascading infections.

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- Stand Alone Cyber Contracts: clauses of repairing actions, through the intervention of a specified cyber security company.
- Prevention or cyber hygiene measures: back-ups, users training, network mapping and inventory etc.

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- Stand Alone Cyber Contracts: clauses of repairing actions, through the intervention of a specified cyber security company.
- Prevention or cyber hygiene measures: back-ups, users training, network mapping and inventory etc.
- ▶ How should we measure the "connectedness" of a given graph G?

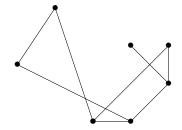
#### Outline of the talk

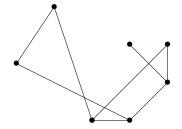
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- Solution and numerical illustrations.

# **Optimization Criteria**

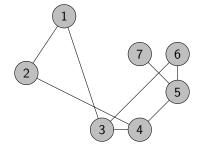
► Given a graph G with adjacency matrix A and degree matrix D, the Laplacian matrix of G is given by L := D – A.

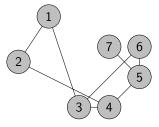
For 
$$x \in \mathbb{R}^n$$
 (*n*= number of nodes),  $x^T L x = \sum_{(i,j) \in E} (x_i - x_j)^2$ .





Choose a numbering of the nodes.





$$L = \begin{pmatrix} 2 & -1 & -1 & 0 & 0 & 0 & 0 \\ -1 & 2 & 0 & -1 & 0 & 0 & 0 \\ -1 & 0 & 3 & -1 & 0 & -1 & 0 \\ 0 & -1 & -1 & 3 & -1 & 0 & 0 \\ 0 & 0 & 0 & -1 & 3 & -1 & -1 \\ 0 & 0 & -1 & 0 & -1 & 2 & 0 \\ 0 & 0 & 0 & 0 & -1 & 0 & 1 \end{pmatrix}$$

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- Algebraic connectivity λ<sub>2</sub>(G)= lowest strictly positive eigenvalue of the Laplacian matrix.
- ▶ Known to describe the "connectedness" of a graph (Fiedler 1973).

# **Optimization Criteria**

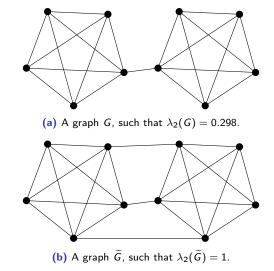
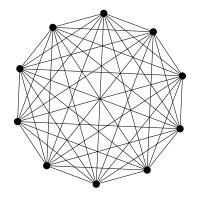


Figure: The lower graph  $\widetilde{G}$  is "more connected" than the upper graph G.

# **Optimization Criteria**



**Figure:** Complete graph  $K_{10}$  with  $\lambda_2(K_{10}) = 10$ .

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- 4. Maximize the algebraic connectivity of L:

where c is a cost function, B a given budget constraint, and A an admissible set of curing rates.

### **Budget constraint**

- 1. Edge *i* interrupted  $\rightarrow \text{loss } Z_i \text{ with } \mathbb{E}[Z_i] = c_i$ , with  $(Z_1, \ldots, Z_m)$  independent and independent of the underlying SIS process.
- 2.  $W(\delta) = \text{total loss on the graph, when the curing rates are given by } \delta \in \mathbb{R}^m$ .
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- **3.**  $\delta_0 = \text{curing rates before any insurance or prevention.}$
- 4. The insurance cost function  $c: \mathbb{R}^m_+ \to \mathbb{R}_+$  is given by

$$c(\delta) := \varphi \left\{ \mathbb{E}[W(\delta_0)] - \mathbb{E}[W(\delta)] \right\}, \tag{1}$$

where  $\varphi : \mathbb{R}_+ \to \mathbb{R}_+$  is non decreasing.

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# **A Solution**

#### Theorem

An optimal solution is given by

$$\delta^{\star}_i = \left(eta \sum_{j=1}^m ilde{a}_{ij} extsf{v}^{\star}_j
ight) rac{1- extsf{v}^{\star}_i}{ extsf{v}^{\star}_i},$$

where  $(1 - v_1^*, \dots, 1 - v_m^*)$  is the *w* component of a solution to the following convex problem in  $\gamma \in \mathbb{R}$ ,  $\mu \in \mathbb{R}$  and  $w \in \mathbb{R}^m$ :

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sup 
$$\gamma$$
  
subject to  $L(w) - \gamma I + \mu 11^t \succcurlyeq 0$   
 $0 \le w \le 1$   
 $\sum_{i=1}^m c_j w_j \le B,$ 

where  $1^t := (1, \ldots, 1) \in \mathbb{R}^n$ , L(w) denotes the Laplacian matrix of the weighted graph  $(V, E_0, w)$ ,  $M \succeq 0$  means that M is positive semidefinite and  $(\tilde{a}_{ij})_{a \leq i,j \leq m}$  are the entries of the edge-adjacency matrix.

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# Idea

Reformulate the problem on  $\delta$  as a problem on v!

#### Lemma 1

 $L = \mathbb{E}[L_{\infty}]$  is the Laplacian matrix associated to the weighted graph  $G = (V, E_0, \overline{v})$ , where each edge  $\ell$  is given the weight  $\overline{v}_{\ell} = 1 - v_{\ell}$ , with  $(v_1, \ldots, v_m)$  satisfying the **steady state equation**.

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Lemma 2

$$\mathbb{E}[\mathcal{W}(\delta_0)] - \mathbb{E}[\mathcal{W}(\delta)] = \sum_{j=1}^m c_j \overline{v}_j.$$

In particular, the cost function c given in (1) is a convex function of v.

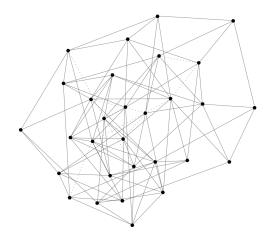
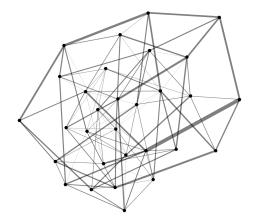


Figure: Erdös-Rényi Random Graph with 102 edges. The dotted edges have optimal  $\delta$  equal to 0.

Solved with the CVX Matlab Package (Boyd, Grant).



**Figure:** Each edge width is proportional to the optimal  $\delta$ .

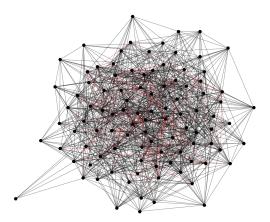
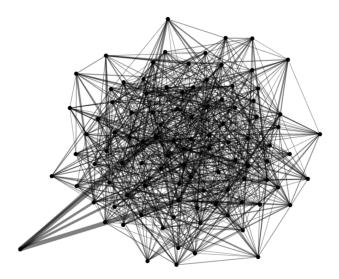
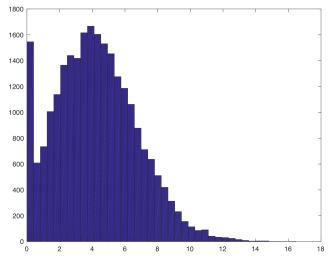


Figure: Erdös-Rényi Random Graph with 1004 edges. Red edges have null  $\delta^{\star}$  values.



**Figure:** Each edge width is proportional to the optimal  $\delta$ .

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**Figure:**  $\delta^*$  values, for a graph with 24978 edges.

### Perspectives

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### Perspectives

- Finite horizon problem and dynamic optimization: stochastic control problem! (Avoids both NIMFA and Steady state approximations)
- Inhomogeneous fast mixing Markov chains (Boyd, Diaconis, Sun, Xiao).
- Simulation and (optimal) importance sampling for Markov chains.
- Optimal stochastic control and games on networks.

# Merci de votre attention